
MAE 552 – Heuristic Optimization

Lecture 25

March 22, 2002

Topic: Tabu Search

A Simple Illustration of Tabu Search

A Simple Version of the short term memory component of the Tabu Search is illustrated in this example.

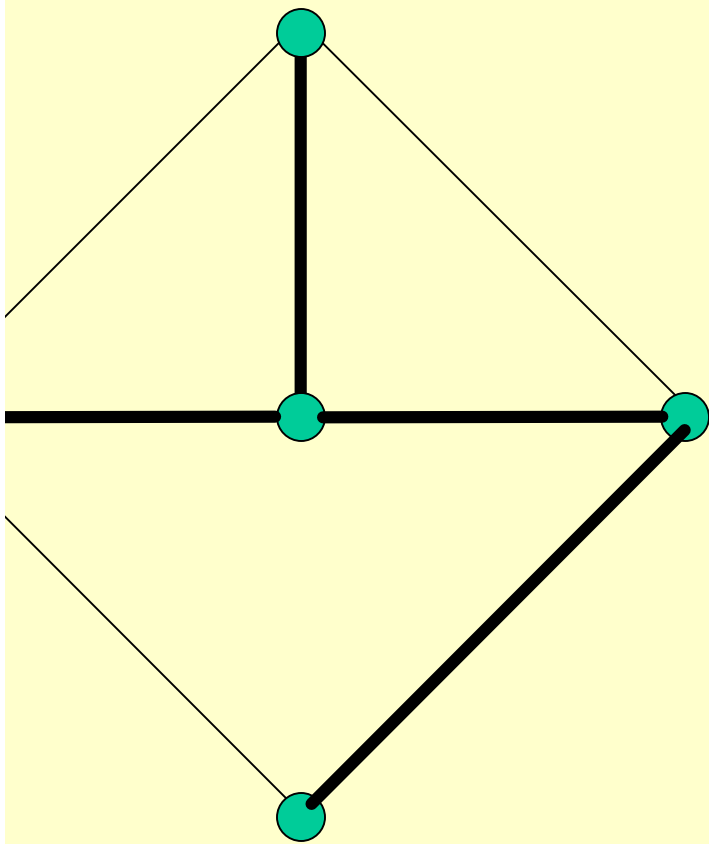
The problem is known as a minimum spanning tree problem

The minimum spanning tree (MST) of a graph defines the sparsest subset of edges that keeps the graph in one connected component.

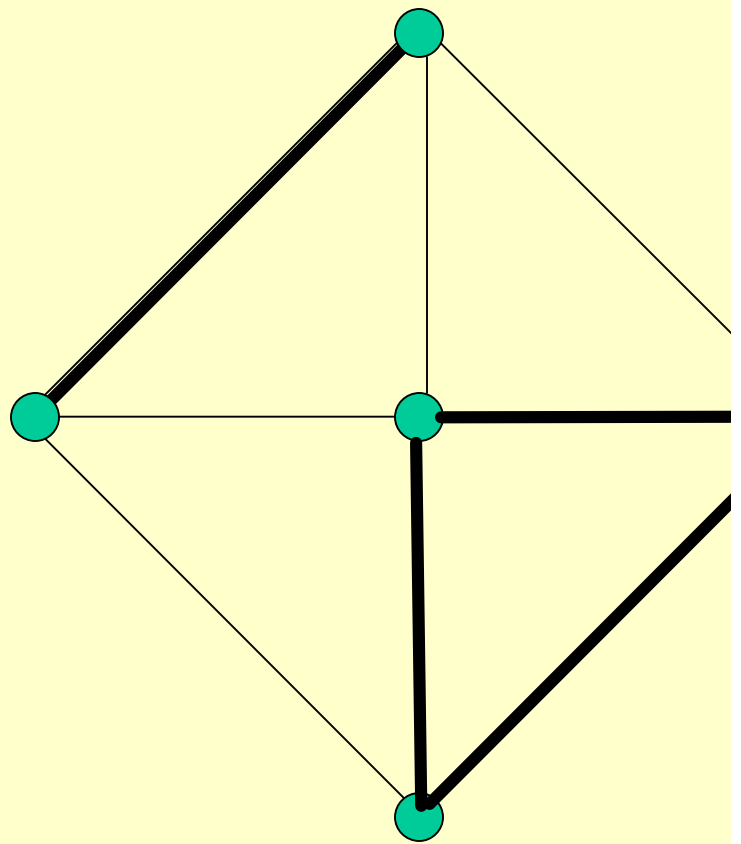
Telephone companies are particularly interested in minimum spanning trees, because the minimum spanning tree of a set of sites defines the wiring scheme that connects the sites using as little wire as possible.

A Simple Illustration of Tabu Search

Legal Spanning Tree

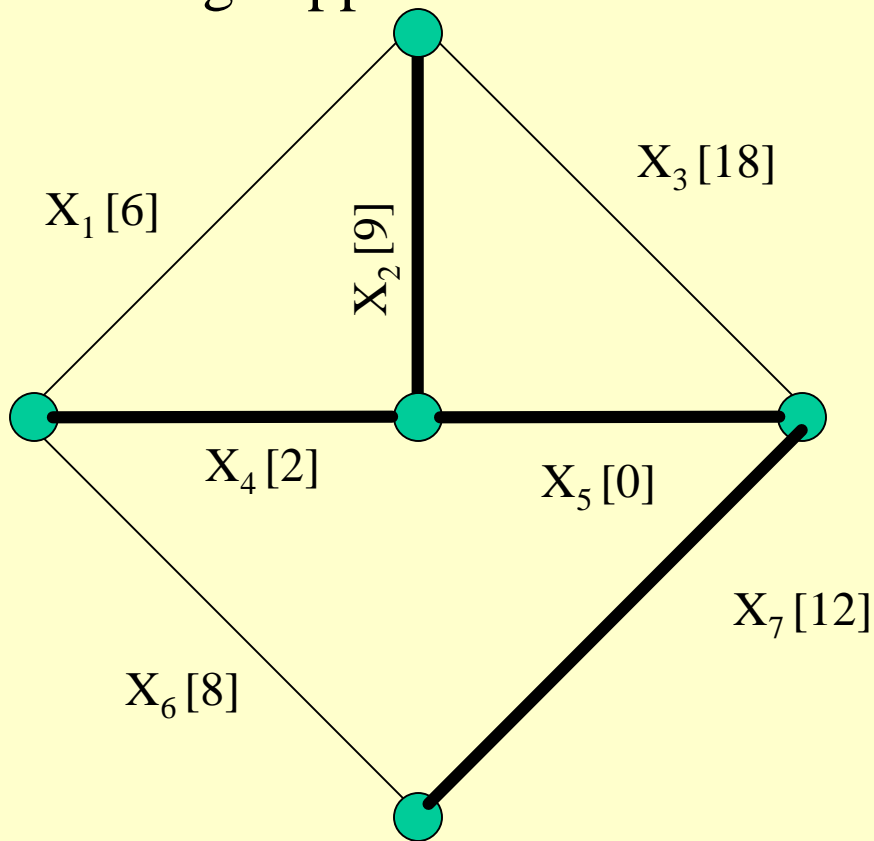


Illegal Spanning Tree



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A solution can be represented in terms of a vector indicating whether or not an edge appears in the solution.



This solution is $(0,1,0,1,1,0,1)$ and $F=23$

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Additionally there are constraints imposed on this problem.

Constraint 1: At most only one of edges 1, 2, or 6 can be used in any tree.
Constraint 2: Edge 1 can be in the tree only if edge 3 is also in the tree.

$$x_1 + x_2 + x_6 \leq 1$$

Constraint 2: Edge 1 can be in the tree only if edge 3 is also in the tree.

$$x_1 \leq x_3$$

To permit the evaluation of the infeasible trees a penalty of 50 is added for each unit violation of a constraint. The a unit violation occurs when the left side of the constraint exceeds the right side by 1.

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To define a Tabu restriction, we have decided to use the *added* edge to be the move attribute assigned Tabu status.

This forbids a future move from dropping the edge as long as it remains Tabu.

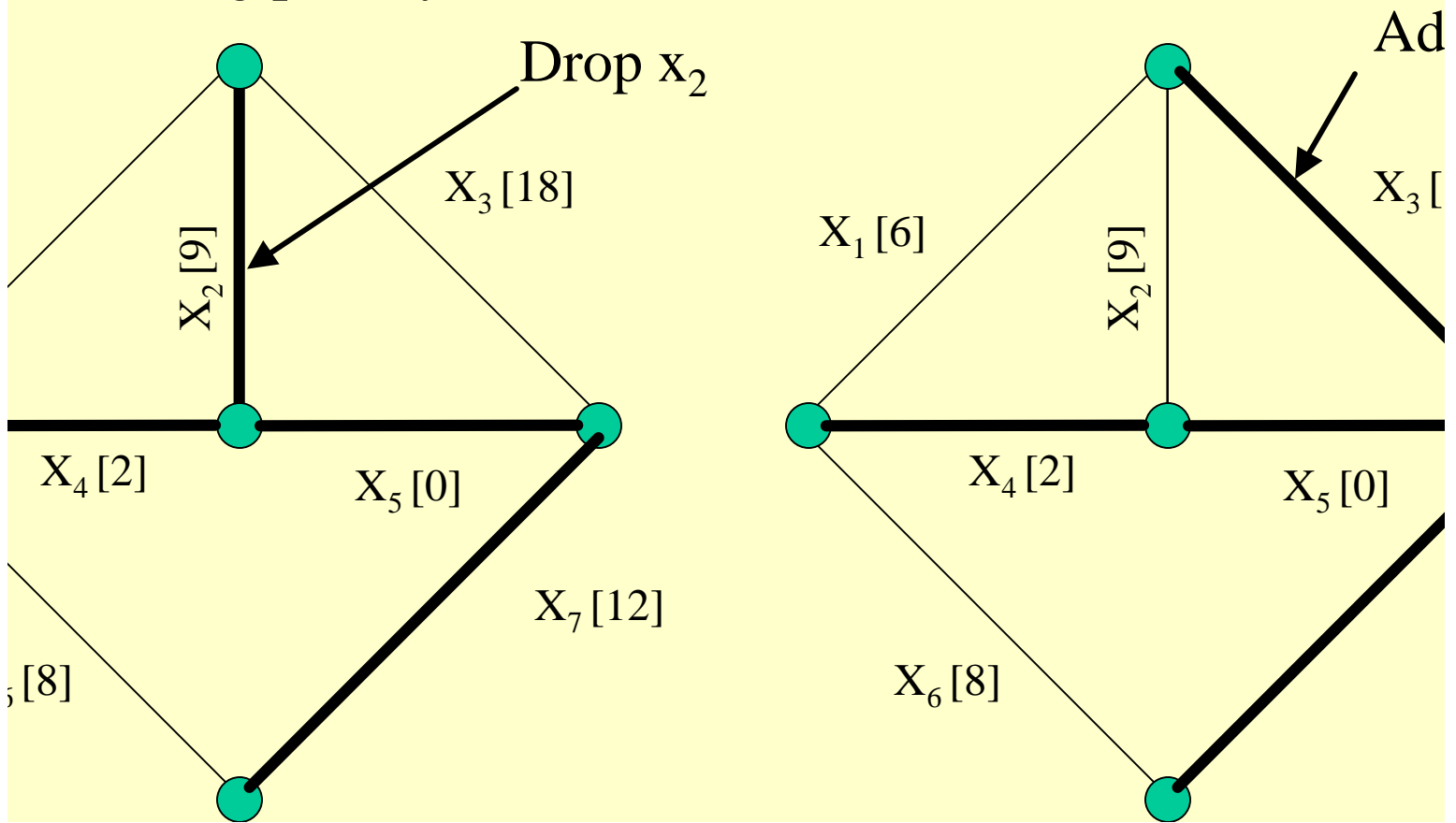
The length of the tabu list for this example is 2.

This move remains Tabu for two iterations and then is dropped from the list.

The aspiration criteria that we have selected is that a tabu restriction can be overridden if the resulting tree is better than that produced so far.

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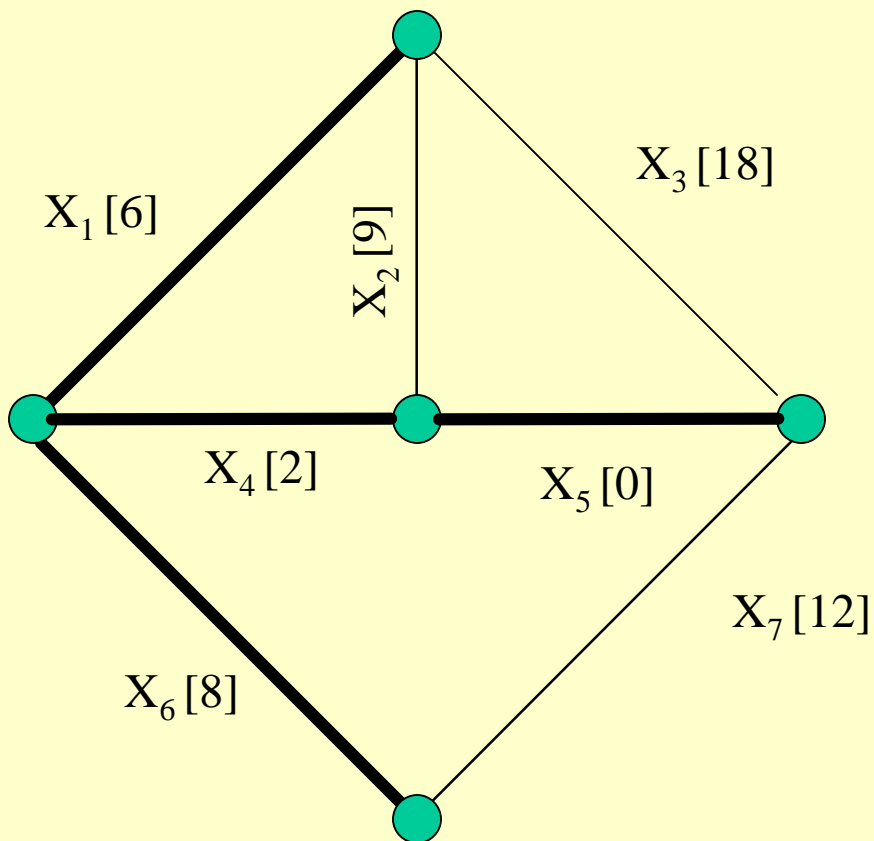
In this example a move will be a standard edge swap that consists of removing an edge and adding an edge to make a new legal tree. The solution selected will be the admissible move with the lowest cost including penalty costs.



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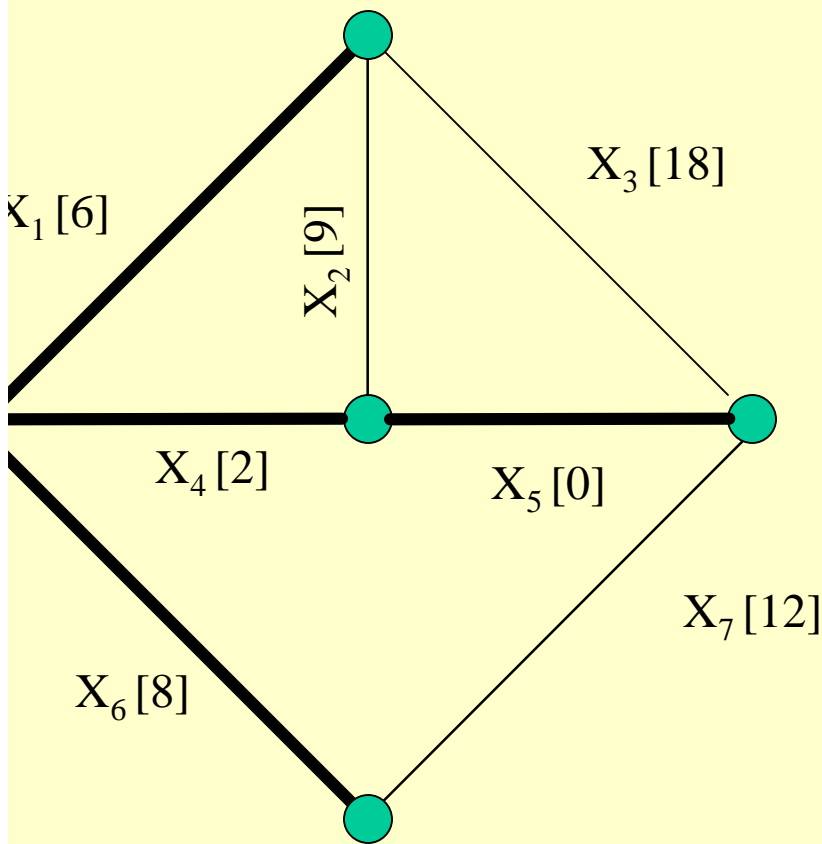
tial Solution Cost = $16 + 100 = 116$

Current Best Point
Infeasible



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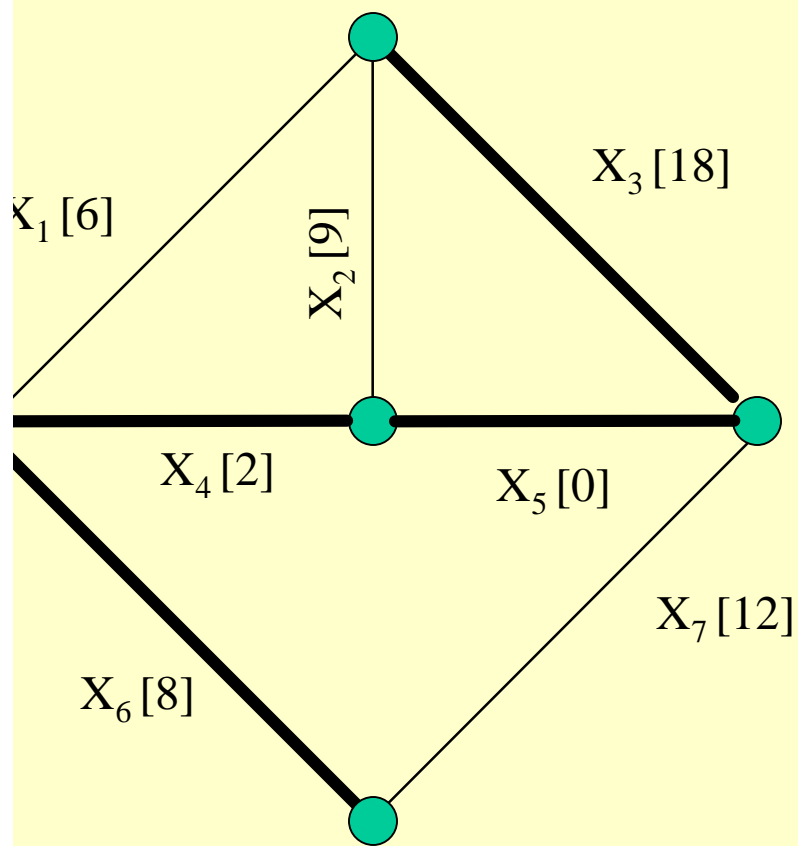
tial Solution Cost = $16 + 100 = 116$



Search neighborhood

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Iteration 2: Current Cost 28 Tabu List: x_3 $M=[0\ 0\ 2\ 0\ 0\ 0\ 0]$

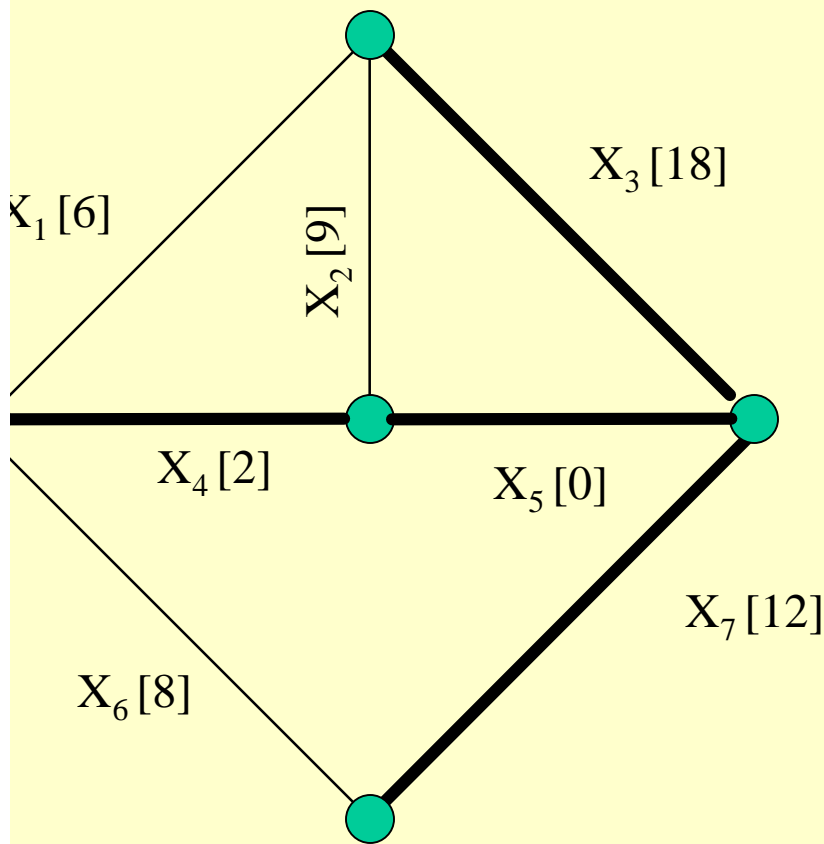


Search neighborhood

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Iteration 3 Current Cost 32 Tabu List: x_3 $M=[0\ 0\ 1\ 0\ 0\ 0\ 2]$

Search neighborhood



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Final Cost 23 Tabu List: x_3 $M=[0\ 2\ 0\ 0\ 0\ 0\ 0]$

